### The Universal $\ell^p$ -Metric on Merge Trees

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CG Week 2022

June 10th, 2022

#### **1** The 5-Minute Overview

### **2** The *p*-Presentation Distance on Merge Trees

**3** Stability and Universality

### Trees for Biology



Fig. 1. Dendrogram drawn based on the matrix of genetic distances among 15 zymodemes of *Trypanosoma cruzi* using UPGMA. The figures on branches indicate the number of times that the branch was observed in 1000 bootstraps. Bootstrap values below 600 are not given. Abbreviations: B, Brazil; Ch, Chile; Co, Colombia; E, Ecuador; G, Guatemala; M, Mexico; Pa, Paraguay; Pe, Peru.

### Trees for Scalar Data



Morozov, Beketayev, and Weber introduced the interleaving distance  $d_I$  on merge trees [4].





N.B.  $d_I(M, N) = d_I(Q, N) = 3$ , but intuitively Q is "closer" to N.

### Cophenetic vectors

- Our vector summaries are subtly different from cophenetic vectors, i.e. the LCA matrix [2, 5, 3], as the length of our vectors is 2n 1 versus  ${}_{n}C_{2} = O(n^{2})$ .
- In particular, the p-cophenetic distance is not Lipschitz stable for  $p \neq \infty$ .





Here  $||f - g||_1 = 2$ , while  $\ell^1$ -cophenetic distance is 3. Instead, we mimic a construction by Bjerkevik and Lesnick [1]. 1 The 5-Minute Overview

### **2** The p-Presentation Distance on Merge Trees

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### Merge Trees as Persistent Sets

A merge tree is a functor  $M \colon \mathbb{R} \to \mathbf{Set}$  that is

- constructible, i.e.  $\exists \tau := \{s_0 < s_1 < \cdots < s_n\} \subset \mathbb{R}$ , such that
  - (i)  $M(s) = \emptyset$  for all  $s < s_0$ , and
  - (ii)  $M(s \le t)$  is an isomorphism whenever  $s, t \in [s_i, s_{i+1})$ , and also for  $s, t \in [s_n, \infty)$ .

• and where |M(t)| = 1 for t sufficiently large.



# Building Blocks for Merge Trees

A strand is a merge tree  $F_s : \mathbb{R} \to \mathbf{Set}$ , for  $s \in \mathbb{R}$ , defined by

$$F_s(t) := \begin{cases} \emptyset & \text{if } t < s, \\ \{*\} & \text{if } t \ge s, \end{cases}$$

with the structure maps all inclusions. We call s birth time of the branch  $F_s$ 



The *p*-Presentation Distance on Merge Trees

## Example Presentation of a Merge Tree

Any merge tree  ${\cal M}$  can be constructed via gluing strands pairwise together.



# Presentation of a Merge Tree

A presentation of a merge tree  $\boldsymbol{M}$  consists of

- generators  $G_i$ 's and relations  $R_j$ 's that are strands;
- together with pairs of underlying merge functions  $f_j, g_j : R_j \to \sqcup_i G_i$  that choose explicit strands for merging.

$$\bigsqcup_{j=1}^{l} R_j \xrightarrow{f} \bigsqcup_{g}^{k} \bigsqcup_{i=1}^{k} G_i \dashrightarrow M,$$

## Presentation Matrix and Label Vector

To a presentation  $P_M$  we have a presentation matrix where

- the *i*-th row corresponds to the *i*-th generator G<sub>i</sub>, labelled by the birth time of G<sub>i</sub>; and
- the *j*-th column corresponds to the *j*-th relation  $R_j$ , labelled by the birth time of  $R_j$ .
- The (i, j)-entry is 1 if  $G_i$  is in the image of  $R_j$  (under f or g) and 0 otherwise.

The label vector  $L(P_M)$  of a  $k \times l$  presentation matrix is the (k + l)-vector where

- $\blacksquare$  the first k entries are the row labels, i.e. heights of leaf nodes, and
- the last l entries are column labels, i.e. the heights of internal nodes.

The *p*-Presentation Distance on Merge Trees



### Compatible Presentations

Two presentations  $P_M$ ,  $P_N$  are **compatible** if their presentation matrices have the same underlying matrix, after forgetting row and column labels.

#### Lemma

Every pair of merge trees M and N, have compatible presentations  $P_M$  and  $P_N$ .

#### Definition

Given  $p \in [1,\infty]$ , the *p*-presentation semi-distance between merge trees M and N is

 $\hat{d}_I^p(M,N) = \inf\{\|L(P_M) - L(P_N)\|_p : P_M \text{ and } P_N \text{ are compatible.}\}$ 

The *p*-Presentation Distance on Merge Trees



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# p-Presentation distance

We see  $\hat{d}_{I}^{p}$  does not satisfy the triangle inequality. Fortunately there is a universal fix.

### Definition

The p-presentation distance between M and N is

$$d_I^p(M,N) := \inf \sum_{i=0}^{n-1} \hat{d}_I^p(Q_i, Q_{i+1}),$$

where we infinize over all finite sequences of merge trees  $M = Q_0, \ldots, Q_n = N$ .

### Theorem (Cardona, C., Lam, Lesnick '21)

- $d_I^{\infty} = d_I$ , i.e., the  $\infty$ -presentation distance equals the interleaving distance.
- For  $p \in [1, \infty]$ ,  $d_I^p$  is a pseudometric.

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## Wasserstein Stability

We extend a lower bound on the interleaving distance due to Morozov et al.

Proposition (CCLL'21)

For  $p \in [1, \infty]$  and merge trees M, N:

 $d^p_{\mathcal{W}}(\mathcal{B}(M), \mathcal{B}(N)) \le d^p_I(M, N).$ 

Here  $d_{\mathcal{W}}^p$  the denotes *p*-Wasserstein distance between barcodes.

# Monotone Cellular Functions

Let X be a finite CW-complex.

- We say  $f: X \to \mathbb{R}$  is monotone if for any face  $\tau$  of  $\sigma$ , one has  $f(\tau) \leq f(\sigma)$ .
- We can define  $||f||_p$  by identifying f with an element of  $\mathbb{R}^{|\mathsf{Cell}(X)|}$ .

#### Theorem (Skraba and Turner, 20')

Let  $f, g: X \to \mathbb{R}$  be monotone cellular functions. Then

 $d^p_{\mathcal{W}}(\mathcal{B}(f), \mathcal{B}(g)) \le ||f - g||_p.$ 

Here  $\mathcal{B}(f)$  is the persistence barcode for the sublevel set filtration of f.

# $\ell^p$ -stability & Universality

We provide an analogue of the interleaving stability for p-presentation distances.

### Theorem ( $\ell^p$ -Stability, CCLL'21)

For any monotone cellular functions  $f, g: X \to \mathbb{R}$ .

 $d_I^p(M_f, M_g) \le ||f - g||_p,$ 

Here  $M_f = \pi_0 \circ S^{\uparrow}(f)$ .

#### Theorem (Universality, CCLL'21)

If d is any distance on merge trees satisfying the above stability property, then  $d \leq d_I^p$ .

# Final Thoughts

- (i) The approach of Bjerkevik and Lesnick seems to generalize to a much broader class of objects. Anything with a notion of presentation where generators and relations have gradings in a metric space should work.
- (ii) However, these metrics feel very complex; NP-most likely.
- (iii) Geometry and stratification theory should guide when the infimum—when passing from the semi-distance to the actual distance—is actually obtained.

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# Thank you for your attention!

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